Assignment 1: Lexing and parsing

Deadline: October 12th, 2022, 23:59

1 Introduction

In this assignment, you will write a program that lexically analyses and parses a piece of text. The grammar used is that of the lambda calculus, as explained during the lectures, and specified below.

The assignment submission must include a program that:

- reads an expression from a file into a character string
- lexically analyzes the character string into a string of tokens
- parses the token string using recursive descent
- outputs a character string in standard format

The program <u>must</u> be able to detect syntax errors and then <u>should</u> report an error. The program <u>must</u> not make use of external libraries. The program <u>should</u> use the least amount of standard library code. The assignment submission <u>must</u> include a Makefile that can be used to compile the program.

The assignment submission must includes a README file that documents:

- The class and group number, and the name(s) of the student(s) who worked on the assignment.
- Whether it is known that the program works correctly, or whether the program has known defects.
- Whether there are any deviations from the assignment, and reasons why.

The README $\underline{\text{may}}$ include an explanation of how the program works, and remarks for improving the assignment. Finally, the assignment submission $\underline{\text{may}}$ include the following two files:

- An archive (positive.tar.gz) of the positive examples used for testing.
- An archive (negative.tar.gz) of the negative examples used for testing.

2 Interface

The program <u>must</u> be compilable (if applicable) and work on the command line. The program can then be invoked using the command line. It accepts one command line argument, namely the file from which it reads.

Input. The program reads a string of characters from the file named by the first command line argument. The program $\underline{\text{may}}$ only work for files which contain only printable ASCII characters and whitespace. The program $\underline{\text{may}}$ accept multiple expressions, one per line. A line is considered terminated by a newline character (\n), or a carriage return and newline character (\n).

Exit status. The program <u>must</u> exit with exit status 0 whenever all expressions in the file are parsed correctly. The program <u>must</u> exist with exit status 1 whenever there is a syntax error, or not enough command line arguments are supplied.

Output. If the program exists with exit status 0 then the program <u>must</u> produce output to the standard output. If the program exists with exit status 1 then an error message <u>may</u> be printed to standard error. The program <u>may</u> print understandable error messages.

3 Grammar

The input file is analyzed using the following Backus-Naur grammar:

$$\langle \exp r \rangle ::= \langle var \rangle \mid ' (' \langle \exp r \rangle ') \mid | ' \langle var \rangle \langle \exp r \rangle \mid \langle \exp r \rangle \langle \exp r \rangle$$

where $\langle \text{var} \rangle$ stands for any variable name. A variable name is alphanumerical: it consists of the letters a-z, A-Z, or the digits 0-9. A variable name <u>must</u> start with a letter from the Latin alphabet, i.e. not with a digit. The grammar <u>should</u> be whitespace insensitive, but whitespace <u>must</u> be used to separate application of two variables. The program <u>may</u> support international variable names (i.e. Unicode), and accept λ instead of λ .

The program <u>must</u> support using parentheses in the input to disambiguate expressions. If no parentheses are used, the order of presedence for the operators is as follows: lambda abstraction groups more strongly than application (i.e. abstraction precedes application), and application associates to the left. The program <u>may</u> support a dot after the lambda abstraction variable, where the dot is parsed in the same way as if an opening parenthesis was inserted with a matching closing parenthesis before the next unmatched closing parenthesis or the end of the expression.

If parsing is succesful, the output of the program <u>must</u> be again acceptable by the program to parse: the program then succesfully parses its own output and <u>should</u> produce the exact same result. The output <u>should</u> be an unambiguous expression, i.e. with sufficiently many parentheses inserted so the parser never applies any of the precedence rules. The output <u>may</u> use the least amount of whitespace and parentheses in its output.

3.1 Positive examples

The following examples are acceptable (the program <u>must</u> work if each expression is put in a separate file, the program <u>should</u> accept multiple expressions each on its own line in a single file):

- (a b)
- abc
- a b c
- a (b c)
- (\ x a b)
- (\x((a) (b)))

3.2 Negative examples

The following examples are not acceptable:

- \ (missing variable after lambda)
- \x (missing expression after lambda abstraction)
- ((x (missing closing parenthesis)
- () (missing expression after opening parenthesis)
- a (b (missing closing parenthesis)
- a (b c)) (input string not fully parsed)

4 Evaluation criteria

The submission will be evaluated on the following criteria:

- Correctness of the program (hard criterium, 60%): is the program correctly implementing the assignment? Are there cases in which the program is implemented incorrectly?
- Readability of the program (soft criterium, 30%): is the program written to be understandable to humans too?
- Efficiency of the program (soft criterium, 10%): is program executing without noticable delay?

In the above text, the words <u>must</u>, <u>should</u>, and <u>may</u> have a special meaning. The assignment is graded with a passing grade if all features that <u>must</u> be implemented are implemented correctly. Higher grades are for submissions that also correctly implement features that <u>should</u> be implemented. Even higher grades are for submissions that also correctly implement features that <u>may</u> be implemented.